

Diffie-Hellman Key Agreement Protocol (DH KAP)

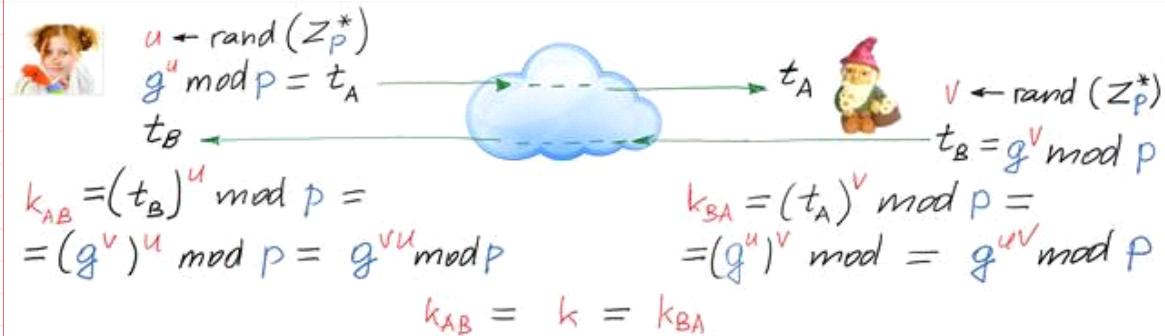
Public Parameters PP=(p, g)

In general, it is a hard problem, but using strong prime p and *Lagrange theorem in group theory* the generator in Z_p^* can be found by random search satisfying two following conditions if p is strongprime: $p = 2*q + 1 \rightarrow q = (p-1)/2$

For all $g \in \Gamma$

$$g^q \neq 1 \pmod{p} \text{ and } g^2 \neq 1 \pmod{p}.$$

```
>> p=genstrongprime(28)
p = 224013599
>> q=(p-1)/2
q = 112006799
>> g=111;
>> mod_exp(g,q,p)
ans = 224013598
>> u=int64(randi(2^28-1))
```



```
>> u=int64(randi(2^28-1))
u = 195162450
>> dec2bin(u)
ans = 1011 1010 0001 1111 0001 0101 0010
>> tA=mod_exp(g,u,p)
tA = 22053505
>> kAB=mod_exp(tB,u,p)
kAB = 196960461
>> v=int64(randi(2^28-1))
v = 212879876
>> tB=mod_exp(g,v,p)
tB = 179573345
>> kBA=mod_exp(tA,v,p)
kBA = 196960461
```

```
>> M='Hello Bob'
M = Hello Bob
>> k=kBA
k = 196960461
>> kh32=dec2hex(k,32)
kh32 = 00000000000000000000000000000000BBD60CD
>> NR=1;
>> fun='e'
fun = e
>> in=M
in = Hello Bob
>> AES128(in,kh32,NR,fun)
new = R H6y N
ans = 18 b3ed521880 cb4836c6de 79810ebe4e
```

% AES128(in,kh32,NR,fun) Advanced Encryption Standard symmetric cipher with key length of **128 bits**
% Encryption is performed for 1 block of length **128 bits** or **16 ASCII symbols**
% in - plaintext/ciphertext of string type: maximum 16 symbols or shorter
% kh32 - shared secret key in hexadecimal number of length=32 (128 bits)
% kh32 can be obtained when shared decimal key **k** is given using commands:
% >> k=int64(randi(2^28))
% k = 160966896
% >> kh32=dec2hex(k,32)
% kh32 = 0000000000000000000000000000000099828F0
%
% NR - Number of Rounds (e.g. NR = 10)
% The smaller NR, the lower security of encryption but the speed of encryption is higher
% The least number of NR is 1 and in this case security lack is evident

```

>> AES128(in,kh32,NR,fun)
new = R H6y N
ans = 18 b3ed521880 cb4836c6de 79810ebe4e
>>
>> fun='d'
fun = d
>> Dh = AES128(Ch,kh32,NR,'d')
Out = 0000000000000048656c6c6f20426f62
Dh = Hello Bob

```

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N >= 10.

```

..... Number of Rounds (e.g. .... 10,
% The smaller NR, the lower security of encryption but the speed of encryption is higher
% The least number of NR is 1 and in this case security lack is evident
%
% fun - letter determining either encryption: fun='e' or decryption: fun='d' functions
%
% Encryption example:
% >> in = 'Hello Bob';
% >> kh32 = '0000000000000000000000000000099828F0';
% >> NR = 10;
% >> Ch = AES128(in,kh32,NR,'e')
% ASCII_e = ?1 ~mV % ciphertext in ASCII format
% Ch = 0f9a2c08d191310fb27ed16d90f45686 % ciphertext in hexadecimal format
%
% Decryption example:
% >> Dh = AES128(Ch,kh32,NR,'d')
% Dh = 00000000000048656c6c6f7720426f62 % decrypted message in hex format
% D = Hello Bob % Decrypted message in ASCII format

```

Public Parameters $PP=(p,g)$

```

>> p
p = 224013599
>> g
g = 111

```

```

PrK = x <-- randi ==> PuK = a = g^x mod p
% input - string array, i.e "asd"
% output - 28 LSBs in hexdecimal form of SHA-256
% >> h=h28('Hello Bob')
h = BD9003E

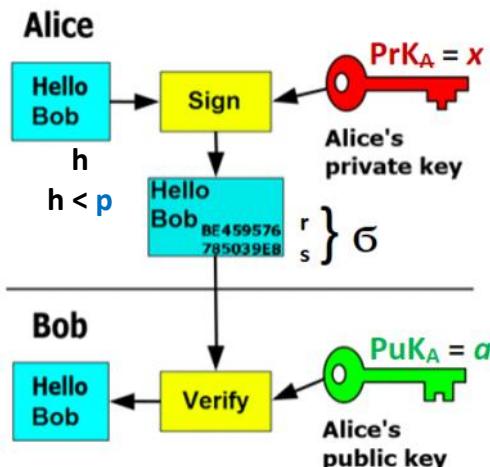
% input - string array, i.e "asd"
% output - 28 LSBs in decimal form of SHA-256
% >> h=hd28('Hello Bob')
h = 198770750

```

```

>> x=int64(randi(2^28-1))
x = 121821633
>> a=mod_exp(g,x,p)
a = 192863004

```



Signature creation for message $M >> p$.

1. Compute decimal h-value $h=H(M)$; $h < p$.
2. Generate $>> i = \text{int64}(\text{randi}(p-1))$ % such that $\text{gcd}(i,p-1)=1$.
3. Compute $i^{-1} \bmod (p-1)$. You can use the function
 $>> i_m1=\text{mulinv}(i, p-1);$
4. Compute $r=g^i \bmod p$.
5. Compute $s=(h \cdot r)x^{-1} \bmod (p-1)$.
6. Signature on h-value h is $\sigma = (r, s)$
 $\text{Sign}(x, h) = \sigma = (r, s).$

```

>> i=int64(randi(p-1))
i = 202651315
>> gcd(i,p-1)
ans = 1
>> i_m1=mulinv(i, p-1)
i_m1 = 59189655
<< mod(i*m1, m1) > 1

```

$$\text{A: } M, \sigma = (r, s) \xrightarrow{\quad} \text{B: }$$

$>> h = \text{hd28}('Hello Bob')$

$>> V1 = \text{mod_exp}(g, h, p)$

```

>> n = na28( Hello Bob )
>> i_m1=mulinv(i, p-1)
i_m1 = 59189655
>> mod(i*i_m1,p-1)
ans = 1
> r=mod_exp(g,i,p)
r = 136585182
>> hmxr=mod(h-x*r,p-1)
hmxr = 76861196
>> s=mod(hmxr*i_m1,p)
s = 106387113

```

```

>> V1 = mod_exp(g,h,p)
>> a_r = mod_exp(a,r,p)
>> r_s = mod_exp(r,s,p)
>> V2 = mod(a_r*r_s,p)

```

2. Signature Verification

A signature $\sigma=(r,s)$ on message M is verified using Public Parameters $PP=(p, g)$ and $PuKA=a$.

1. Bob computes $h=H(M)$.
2. Bob verifies if $1 < r < p-1$ and $1 < s < p-1$.
3. Bob calculates $V1=g^h \text{ mod } p$ and $V2=a^r r^s \text{ mod } p$, and verifies if $V1=V2$.

The verifier Bob accepts a signature if all **conditions** are satisfied during the signature creation and **rejects** it otherwise.